

# Handout Chapter I section 11

## Chapter II section 1-2

Seminar Constructible Sets

Olle Torstensson & Tristan van der Vlugt

**Reminder** (lemma I.9.11). Let  $\Phi(\vec{x})$  be a  $\Sigma_0$  formula of LST and let  $\varphi(\vec{x})$  be its counterpart in  $\mathcal{L}$ , then

$$\text{ZF} \vdash \forall u (\forall \vec{x} \in u) \left[ \Phi^u(\vec{x}) \leftrightarrow \vDash_u \varphi(\vec{x}) \right]$$

**Reminder** (lemma I.9.15). Let  $\Phi(\vec{x})$  be a  $\Sigma_0$  formula of LST and let  $\varphi(\vec{x})$  be its counterpart in  $\mathcal{L}$ , then

$$\text{ZF} \vdash \text{“For } M \text{ a transitive set, } (\forall \vec{x} \in M) \left[ \Phi(\vec{x}) \leftrightarrow \vDash_M \varphi(\vec{x}) \right] \text{”}$$

### 1.11) Kripke-Platek Set Theory & Admissible sets

**Definition 1.** KP (Kripke-Platek set theory) is the theory given by the axioms of BS enriched with  $\Delta_0$ -collection:

- (i) Axiom of Extensionality
- (ii) Axiom of Induction
- (iii) Axiom of Pairing
- (iv) Axiom of Union
- (v) Axiom of Infinity
- (vi) Axiom of Cartesian Product
- (vii) Axiom schema of  $\Delta_0$ -Comprehension
- (viii) Axiom schema of  $\Delta_0$ -Collection

**Reminder** (amenable sets). A set  $M$  is **amenable**, if it is transitive and satisfies

- (i)  $(\forall x \in M)(\forall y \in M)(\{x, y\} \in M)$
- (ii)  $(\forall x \in M)(\bigcup x \in M)$
- (iii)  $\omega \in M$
- (iv)  $(\forall x \in M)(\forall y \in M)(x \times y \in M)$
- (v) If  $R \subseteq M$  is  $\Sigma_0(M)$ , then  $(\forall x \in M)(R \cap x \in M)$

**Definition 2.** A set  $M$  is **admissible**, if  $M$  is amenable and satisfies

- (vi) If a relation  $R \subseteq M \times M$  is  $\Sigma_0(M)$  and  $(\forall x \in M)(\exists y \in M)(y R x)$ , then for any  $u \in M$  there is a  $v \in M$  such that  $(\forall x \in u)(\exists y \in v)(y R x)$

**Lemma 3.** Let  $\Phi(x, y, \vec{p})$  be  $\Sigma_1$ , and  $\Psi(z, \vec{p})$  be  $\Delta_1^{\text{KP}}$  then

$$\text{KP} \vdash \forall \vec{p} [\forall x \exists y \Phi(x, y, \vec{p}) \rightarrow \forall u \exists v (\forall x \in u)(\exists y \in v) \Phi(x, y, \vec{p})] \quad (\Sigma_1\text{-Collection})$$

$$\text{KP} \vdash \forall \vec{p} [\forall x \exists y \forall z (z \in y \leftrightarrow z \in x \wedge \Psi(z, \vec{p}))] \quad (\Delta_1\text{-Comprehension})$$

**Lemma 4** (Recursion Theorem). Let  $G$  be a total  $(n+2)$ -ary  $\Sigma_1^{\text{KP}}$  function over  $V$ . Then there is a total  $(n+1)$ -ary  $\Sigma_1^{\text{KP}}$  function  $F$  over  $V$  such that

$$\text{KP} \vdash F(\vec{x}, y) = G(\vec{x}, y, (F(\vec{x}, z) \mid z \in y)).$$

## 2.1) Definition of the Constructible Universe

**Reminder** (definition of  $x$ -definable). A subset  $y$  of  $x$  is  $x$ -**definable** iff there is a formula  $\varphi(v)$  of  $\mathcal{L}_x$  such that  $y = \{a \in x \mid \models_x \varphi(\hat{a})\}$ . Define  $\text{Def}(x)$  to be the set of  $x$ -definable subsets of  $x$ .

$\text{Def}(u)$  can be defined in LST as:

$$\begin{aligned} v = \text{Def}(u) &\leftrightarrow (\forall x \in v)(\exists \varphi) \left[ \text{Fml}(\varphi, u) \wedge \text{Fr}(\varphi, \{v_0\}) \right. \\ &\quad \wedge x = \left. \{z \in u \mid (\exists \psi)(\text{Sub}(\psi, \varphi, v_0, \hat{z}) \wedge \text{Sat}(u, \psi))\} \right] \\ &\quad \wedge (\forall \varphi) \left[ (\text{Fml}(\varphi, u) \wedge \text{Fr}(\varphi, \{v_0\})) \right. \\ &\quad \left. \rightarrow (\exists x \in v)(x = \{z \in u \mid (\exists \psi)(\text{Sub}(\psi, \varphi, v_0, \hat{z}) \wedge \text{Sat}(u, \psi))\}) \right] \end{aligned}$$

**Definition 5.** By recursion on  $\alpha \in \text{On}$  we define

$$L_0 = \emptyset, \quad L_{\alpha+1} = \text{Def}(L_\alpha), \quad L_\gamma = \bigcup_{\alpha < \gamma} L_\alpha \text{ if } \gamma \text{ is limit}, \quad L = \bigcup_{\alpha \in \text{On}} L_\alpha.$$

The set  $L$  is the **constructible universe**. A set  $x$  is **constructible** iff  $x \in L$ .

**Lemma 6.** The constructible hierarchy has the following properties:

- (i)  $\alpha \leq \beta$  implies  $L_\alpha \subseteq L_\beta$ .
- (ii)  $L_\alpha$  is transitive for each  $\alpha$ . ( $L$  is transitive).
- (iii)  $L_\alpha \subseteq V_\alpha$  for all  $\alpha$ .
- (iv)  $\alpha < \beta$  implies  $\alpha \in L_\beta$  and  $L_\alpha \in L_\beta$ . ( $\text{On} \subseteq L$ ).
- (v)  $L \cap \alpha = L_\alpha \cap \text{On} = \alpha$  for all  $\alpha$ .
- (vi)  $L_\alpha = V_\alpha$  for  $\alpha \leq \omega$ .
- (vii)  $|L_\alpha| = |\alpha|$  for  $\alpha \geq \omega$ .

**Definition 7.** Let  $M$  be a transitive proper class and  $T$  a theory in LST, then  $M$  is an **inner model** of  $T$  if for every axiom  $\Phi$  of  $T$  we have  $T \vdash \Phi^M$ .

**Theorem 8.** The class  $L$  is an inner model of ZF. In fact  $\text{ZF} \vdash (\text{AoC})^L$  (this will be treated next week).

## 2.2) Constructibility

**Lemma 9.**  $L_\alpha$  is amenable for each uncountable limit ordinal  $\alpha$ .

**Reminder** (uniformly  $\Delta_n^M$ ). Given a family of classes  $\mathcal{F}$  and a class  $\mathcal{A} = \{\vec{x} \mid \Phi(\vec{x})\}$  of  $m$ -tuples defined by LST formula  $\Phi$ , we say  $\mathcal{A}$  is uniformly  $\Delta_n^M$  if there are a  $\Sigma_n$  formula  $\varphi_1(\vec{x})$  and  $\Pi_n$  formula  $\varphi_2(\vec{x})$  of  $\mathcal{L}$  such that for each  $\mathbf{M} \in \mathcal{F}$  we have  $A \cap M^m = \{\vec{x} \mid \models_{\mathbf{M}} \varphi_i(\vec{x})\}$ .

**Remark 10.** Devlin states that the class defined by  $\text{Sat}$  (that is  $\{(u, \varphi) \mid \text{Sat}(u, \varphi)\}$ ) is uniformly  $\Delta_1^M$  for amenable sets  $M$ . This is **false**, as Mathias gives a countermodel. It can be fixed by defining  $\mathcal{S}(x)$  to be the set of finite subsets of  $x$ , and letting  $M$  be an  $\mathcal{S}$ -amenable set, which is an amenable set such that  $x \in M$  implies  $\mathcal{S}(x) \in M$ . Now the class  $\text{Sat}$  is  $\Delta_1^{M^{\mathcal{S}}}$  for  $\mathcal{S}$ -amenable sets  $M^{\mathcal{S}}$ .

**Definition 11.** For an LST formula  $\Phi(\vec{x})$ , let  $(*)$  and  $(**)$  be the following properties:

- $(*)$   $\Phi(\vec{x})$  is  $\Delta_1^{\text{KP}}$ .

(\*\*) The class  $\{\vec{x} \mid \Phi(\vec{x})\}$  is uniformly  $\Delta_1^{L_\alpha}$ , for all uncountable limit ordinals  $\alpha$ .

We define the following LST formulas:

$$\begin{aligned} \text{Seq}(y, x) : & \quad \exists f \left[ \text{“}f \text{ is a function”} \wedge \text{dom}(f) = \omega \wedge f(0) = \emptyset \wedge y = \bigcup \text{ran}(f) \right. \\ & \quad \wedge (\forall n \in \omega)(\forall s \in f(n+1))(\exists t \in f(n))(\exists a \in x) \left( s = t \cup \{(a, n)\} \right) \\ & \quad \left. \wedge (\forall n \in \omega)(\forall s \in f(n))(\forall a \in x)(\exists t \in f(n+1)) \left( t = s \cup \{(a, n)\} \right) \right] \\ \text{Pow}(y, x) : & \quad \exists z \left[ \text{Seq}(z, x) \wedge y = \{\text{ran}(u) \mid u \in z\} \right] \end{aligned}$$

**Lemma 12.**  $\text{Seq}(y, x)$  satisfies (\*) and (\*\*).

**Lemma 13.**  $\text{Pow}(y, x)$  satisfies (\*) and (\*\*).

**Lemma 14.** There is an LST formula  $D(v, u)$  such that  $D(v, u) \leftrightarrow v = \text{Def}(u)$ . Moreover,  $D(v, u)$  satisfies (\*) and (\*\*).

**Lemma 15.** There is an LST function  $G(f, \alpha)$  such that  $G$  is true iff  $f : \alpha + 1 \rightarrow L$  sends  $\beta \mapsto L_\beta$ . Moreover,  $G(f, \alpha)$  satisfies (\*) and (\*\*).

Let  $H(x, \alpha)$  be the LST formula that says “ $x = L_\alpha$ ”, which is defined as  $H(x, \alpha) \leftrightarrow \exists f[G(f, \alpha) \wedge (x = f(\alpha))]$ .

**Lemma 16.**  $H(x, \alpha)$  satisfies (\*) and (\*\*).

**Lemma 17.** If  $M$  is an inner model of KP, and  $\alpha$  an ordinal, then  $L_\alpha \in M$  and  $(L_\alpha)^M = L_\alpha$ . Hence  $(L)^M = L$ .

**Corollary 18.**

- (i) If  $M$  is an admissible set and  $\lambda = \sup(M \cap \text{On})$ , then for  $\alpha < \lambda$  we have  $(L_\alpha)^M = L_\alpha$ . Hence  $(L)^M = L_\lambda$ .
- (ii) For any  $\alpha$  we have  $(L_\alpha)^L = L_\alpha$ . Hence  $(L)^L = L$ .
- (iii) For  $\alpha > \omega$  a limit ordinal and  $\gamma < \alpha$  we have  $(L_\gamma)^{L_\alpha} = L_\gamma$ . Hence  $(L)^{L_\alpha} = L_\alpha$ .

**Lemma 19.** The LST formula “ $x \in L$ ” (that is,  $x$  is constructible) is  $\Sigma_1^{\text{KP}}$ .

**Definition 20.** The **Axiom of Constructibility** (also *Hypothesis of Constructibility*) is the statement  $\forall x(x \in L)$ . Equivalently this means  $V = L$ .

**Theorem 21.**  $\text{ZF} \vdash (V = L)^L$ .

## 1 Exercises

**Exercise 1.** Prove that the Axiom of Cartesian Product holds in the inner model  $L$

$$\text{ZF} \vdash (\forall x \forall y \exists z \forall u [u \in z \leftrightarrow (\exists a \in x)(\exists b \in y)(u = \langle a, b \rangle)])^L$$

without assuming II.1.2. However, you may assume that

$$\text{ZF} \vdash \forall x \forall y \exists z \forall u [u \in z \leftrightarrow (\exists a \in x)(\exists b \in y)(u = \langle a, b \rangle)].$$

**Exercise 2.** Let  $\text{Add}(s, a, b)$  be the LST formula that states “ $s = a + b$ ”  $\wedge$  “ $a, b, s$  are natural numbers”.

- a) Prove that  $\text{Add}(s, a, b)$  is  $\Delta_1^{\text{KP}}$ . (\*)
- b) Prove that the class defined by  $\text{Add}$  is uniformly  $\Delta_1^{L^\alpha}$  for  $\alpha > \omega$  limit. (\*\*)