

# Seminar on Set Theory

Hand-in exercise 15

January 15, 2016

## Exercise 1.

- (i) First we show that Collection implies Replacement. If  $\forall x \in u \exists! y \varphi(x, y)$  is the case, then by Collection we can form a set  $v$  such that  $\forall x \in u \exists y \in v \varphi(x, y)$ . We apply Separation on  $v$  to find the set  $w = \{y \in v : \exists x \in u \varphi(x, y)\}$ . Clearly if  $y \in w$  then  $\exists x \in u \varphi(x, y)$ ; on the other hand, if for some  $y$  we have  $\exists x \in u \varphi(x, y)$ , then because of uniqueness of  $y$  we must have  $y \in v$  and so  $y \in w$ .

Now to show that Replacement implies Collection. Suppose that  $\forall x \in u \exists y \varphi(x, y)$ . Define  $\psi(x, w)$  by  $\exists \alpha \in \text{ORD}[w = V_\alpha \wedge \forall \beta \in \alpha[\alpha = \beta \leftrightarrow \exists y \in V_\beta \varphi(x, y)]]$ . By Regularity we have  $V = \bigcup_{\alpha \in \text{ORD}} V_\alpha$ , hence  $\forall x \in u \exists! w \psi(x, w)$ . By Replacement, there is some  $u$  such that  $\forall w[w \in u \leftrightarrow \exists x \in u \psi(x, w)]$ . By Union we can form  $\bigcup u$ , of which we recognize that  $\forall x \in u \exists y \in \bigcup u \varphi(x, y)$ , which concludes the proof.  $\square$

*Points awarded:*  $\frac{1}{2}$  point per correctly proven direction of the equivalence, where the solution is required to explicitly mention the crucial axioms.

- (ii) First we show that Set Induction implies Regularity. From Set Induction we can conclude that  $V = \bigcup_{\alpha \in \text{ORD}} V_\alpha$ , so we can use the notion of rank of a set. Thus if  $u \neq \emptyset$  is any set, then since the ordinals are well-founded it must have an element  $x \in u$  of lowest rank, for which it must hold that  $x \cap u = \emptyset$  as required. [There are more direct and elementary proofs as well.]

Now to show that Regularity implies Set Induction. Suppose  $\forall x[\forall y \in x \varphi(y) \rightarrow \varphi(x)]$ , so that for any  $x$  we have  $\neg \varphi(x) \rightarrow x^- \neq \emptyset$ , where  $x^-$  denotes the set  $\{y \in x : \neg \varphi(y)\}$ . Assume we have some set  $x$  such that  $\neg \varphi(x)$ . Then there is some  $z_1 \in x^-$ : but then since  $\neg \varphi(z_1)$  again holds, there is in turn some  $z_2 \in z_1^-$ , which by iteration results in an infinite sequence  $x \ni z_1 \ni z_2 \ni \dots$ . By Replacement these form a set which contradicts Regularity, so we must have  $\forall x \varphi(x)$ , which completes the proof.  $\square$

*Points awarded:* 1 point per correctly proven direction of the equivalence, where the solution is required to explicitly mention the crucial axioms.

- (iii) First we show that Fullness implies Subset Collection. For  $u$  and  $v$ , let  $F_{u,v}$  be a full set of total relations between  $u$  and  $v$  as given by Fullness. For any relation  $R$  between  $u$  and  $v$  we can form by Restricted Separation the set  $v_R = \{y \in v : \exists x \in u \langle x, y \rangle \in R\}$ . By Strong Collection there is a set  $W_{u,v} = \{v_R : R \in F_{u,v}\}$ . Now let  $z$  be arbitrary, and suppose that  $\forall x \in u \exists y \in v \varphi(x, y, z)$ . We would like some total relation  $R_z$  between  $u$  and  $v$  which is based on  $z$ . We cannot straightforwardly define  $R_z$  as  $\{\langle x, y \rangle : x \in u \wedge y \in v \wedge \varphi(x, y, z)\}$  since  $\varphi(x, y, z)$  is generally not restricted. Therefore we instead obtain  $R_z$  via Strong Collection on the formula  $\psi(x, p)$  which is  $\exists y \in v[p = \langle x, y \rangle \wedge \varphi(x, y, z)]$ . We see that  $R_z$  thus defined is indeed a total relation between  $u$  and  $v$ , and that  $\varphi(x, y, z)$  is the case whenever  $\langle x, y \rangle \in R_z$ . Now there is some  $R \in F_{u,v}$  such that  $R \subseteq R_z$ , along with corresponding  $v_R \in W_{u,v}$ . Since  $R$  is again total, for all  $x \in u$  there is some  $y \in v_R$  such that  $\langle x, y \rangle \in R$ , but then also  $\langle x, y \rangle \in R_z$  and so  $\varphi(x, y, z)$  holds. On the other hand, if  $y \in v_R$ , then by definition there is some  $x \in u$  such that  $\langle x, y \rangle \in R$  and so  $\langle x, y \rangle \in R_z$ , by which we again obtain  $\varphi(x, y, z)$ . This shows that Subset Collection

holds with  $W_{u,v}$  as witness.

Now to show that Subset Collection implies Fullness. For  $u$  and  $v$ , consider the formula  $\varphi(x, p, z)$  given by  $p \in z \wedge \exists y \in v (p = \langle x, y \rangle)$ . By Subset Collection we find a set  $w_{u, u \times v}$  such that for any  $z$  we have  $\forall x \exists p \in u \times v \varphi(x, p, z) \rightarrow \exists P_z \in w_{u, u \times v} [\forall x \in u \exists p \in P_z \varphi(x, p, z) \wedge \forall p \in P_z \exists x \in u \varphi(x, p, z)]$ . Now let  $z$  be any total relation between  $u$  and  $v$ , so that for every  $x \in u$  there is some  $y \in v$  such that  $\langle x, y \rangle \in z$ . Then clearly  $\forall x \exists p \in u \times v \varphi(x, p, z)$  is the case, hence there is some  $P_z \in w_{u, u \times v}$  such that  $\forall x \in u \exists p \in P_z \varphi(x, p, z)$  as well as  $\forall p \in P_z \exists x \in u \varphi(x, p, z)$ . The latter tells us that  $P_z \subseteq z$  is a relation between  $u$  and  $v$ , and the former gives us moreover that  $P_z$  is total, thus for any total relation  $z$  the set  $w_{u, u \times v}$  contains a total relation  $P_z$  which refines it. Since being a total relation is restricted, by Restricted Separation we may consider  $f_{u, u \times v} = \{R \in w_{u, u \times v} : \text{TRel}(R, u, v)\}$ , which then witnesses Fullness for our arbitrary  $u$  and  $v$ .  $\square$

*Points awarded:* 1 point per correctly proven direction of the equivalence, where the solution is required to explicitly mention the crucial axioms.

### Exercise 2.

- (i) If  $\alpha < \beta^+$ , then  $\alpha \in \beta \cup \{\beta\}$ , so  $\alpha \in \beta$  or  $\alpha = \beta$ . In the first case, the transitivity of  $\beta$  yields that  $\alpha \subset \beta$ . In the second case, we also have  $\alpha \subset \beta$ . We conclude that  $\alpha \leq \beta$ .  $\square$
- (ii) Suppose that  $\forall \alpha, \beta (\alpha \leq \beta \rightarrow \alpha < \beta^+)$  holds. We let  $\alpha = 0$  and  $\beta = \{0 \mid \phi\}$ . Then clearly  $0 \subset \{0 \mid \phi\}$ , so we get  $0 \in \{0 \mid \phi\}^+$ . That is,  $0 \in \{0 \mid \phi\}$  or  $0 = \{0 \mid \phi\}$ . In the first case, it follows that  $\phi$ , while in the second case, it follows that  $\neg\phi$ .  $\square$   
*Alternatively, one may take  $\alpha = \{0 \mid \phi\}$  and  $\beta = 1$ .*
- (iii) Suppose that  $\forall \alpha, \beta, \gamma (\alpha \leq \beta < \gamma \rightarrow \alpha < \gamma)$  holds. We let  $\gamma = \beta^+$ . Since  $\beta < \beta^+$  always holds, we now get  $\forall \alpha, \beta (\alpha \leq \beta \rightarrow \alpha < \beta^+)$ , which implies LEM, by exercise (ii).  $\square$
- (iv) We apply this to  $\alpha = \{0 \mid \phi\}$ . Suppose that  $\{0 \mid \phi\}$  is a weak limit. Suppose we have a  $\beta \in \{0 \mid \phi\}$ , then there must also be a  $\gamma \in \{0 \mid \phi\}$  such that  $\beta \in \gamma$ . But since  $\beta$  and  $\gamma$  are both in  $\{0 \mid \phi\}$ , we have  $\beta = \gamma = 0$ . However, we clearly do not have  $0 \in 0$ , contradiction. This shows that  $\neg \exists \beta \in \{0 \mid \phi\}$ , so  $\{0 \mid \phi\} = 0$ , whence  $\neg\phi$  holds. If  $\{0 \mid \phi\} = \beta^+$  for some ordinal  $\beta$ , then we have  $\beta \in \{0 \mid \phi\}$ , so  $\beta = 0$ . This means that  $\{0 \mid \phi\} = 0^+ = \{0\}$ , so  $\phi$  holds.  $\square$

*One point was awarded for correctly handling the case in which  $\{0 \mid \phi\}$  is a successor, and one point for the case in which it is a weak limit.*