

# Hand-in 16: Seminar Set Theory

January 25, 2016

## Problem 1: Category Theory

When defining products, exponential objects, and monomorphisms during the lecture, we first considered the construction in terms of morphisms  $1 \rightarrow A$ , and then generalised this to morphisms  $Z \rightarrow A$  for any object  $Z$ .

For any one of these constructions:

(a, 0.5 points): Show that in *Set*, the definition remains equivalent if we replace “for any object  $Z$ ” with “for some terminal object  $1$ ”.

(b, 0.5 points): Show that this is not necessarily the case for an arbitrary topos. (Hint: *Set*<sup>2</sup> with all operations defined componentwise is a topos.)

Now let  $\mathcal{C}$  be an arbitrary category with a terminal object  $1$ .

(c, 0.5 point): Show that for any object  $A$  of  $\mathcal{C}$ , any morphism  $x : 1 \rightarrow A$  is a monomorphism.

Assume furthermore that  $\mathcal{C}$  is an elementary topos.

(d, 0.5 point): Let  $0$  be an object such that for all  $Z \in \text{Ob}$ , there exists exactly one morphism  $f : 0 \rightarrow Z$ . Show that  $f$  is a monomorphism.

Let  $A \in \text{Ob}$  and suppose the product  $A \times 1$  exists in  $\mathcal{C}$ . We say two objects  $A$  and  $B$  are isomorphic if there exist morphisms  $f : A \rightarrow B$  and  $g : B \rightarrow A$  such that  $g \circ f = \text{id}_A$  and  $f \circ g = \text{id}_B$ .

(e, 0.5 point): Show that  $A$  and  $A \times 1$  are isomorphic.

## Problem 2: Truth in *Set*<sub>H</sub>

Let  $H$  be a complete Heyting algebra.

**Definition 1.** An  $H$ -set is a pair  $(X, \delta)$  where  $\delta : X \times X \rightarrow H$  satisfying

$$\begin{aligned} \delta(x, x') &= \delta(x', x) & \forall x, x' \in X \\ \delta(x, x') \wedge \delta(x', x'') &\leq \delta(x, x'') & \forall x, x', x'' \in X. \end{aligned}$$

**Definition 2.** An  $H$ -valued functional relation (for brevity:  $H$ -function) from  $(X, \delta_X)$  to  $(Y, \delta_Y)$  is a function  $f : X \times Y \rightarrow H$  satisfying

$$\begin{aligned} \delta_X(x, x') \wedge f(x', y) &\leq f(x, y) & \forall x, x' \in X, y \in Y \\ f(x, y) \wedge \delta_Y(y, y') &\leq f(x, y') & \forall x \in X, y, y' \in Y \\ f(x, y) \wedge f(x, y') &\leq \delta_Y(y, y') & \forall x \in X, y, y' \in Y \\ \bigvee_{y \in Y} f(x, y) &= \delta_X(x, x) & \forall x \in X. \end{aligned}$$

Define composition of  $f : (X, \delta_X) \rightarrow (Y, \delta_Y)$  and  $g : (Y, \delta_Y) \rightarrow (Z, \delta_Z)$  by

$$(g \circ f)(x, z) = \bigvee_{y \in Y} (f(x, y) \wedge g(y, z)).$$

(2.5 points): We claim  $Set_H$  is a topos. Find the terminal object  $1$ , the truth-value object  $\Omega$ , and the  $true : 1 \rightarrow \Omega$  morphism and show that they satisfy definitions 2 and 7 of the handout.