

**Exam Topos Theory, June 16, 2022, 14:00–17:00**

This exam consists of 4 exercises. Every exercise is worth 10 points; if an exercise consists of more than one part, it is indicated what each part is worth. The grade  $W$  for the written exam is your total number of points divided by 4. Your final grade is the maximum of  $W$  and  $\frac{7W+3H}{10}$  where  $H$  is your result from the homework exercises.

Advice: first do those exercises for which you see a solution right away. Then start thinking about the harder ones. Good luck!

**Exercise 1** First, let us establish some terminology and notation. For an object  $A$  of a topos  $\mathcal{E}$ , the monomorphism  $\in_A \rightarrow \Omega^A \times A$  is the subobject classified by the evaluation map  $\Omega^A \times A \rightarrow \Omega$ . For an object  $C$ , the map  $\{\cdot\} : C \rightarrow \Omega^C$  is the exponential transpose of the map  $\Delta : C \times C \rightarrow \Omega$  which classifies the diagonal  $\delta : C \rightarrow C \times C$ . For a subobject  $R$  of  $B \times A$  we say that a map  $r : B \rightarrow \Omega^A$  names  $R$  (or, is a name for  $R$ ), if there is a pullback square

$$\begin{array}{ccc} R & \longrightarrow & \in_A \\ \downarrow & & \downarrow \\ B \times A & \xrightarrow{r \times \text{id}_A} & \Omega^A \times A \end{array}$$

It can be used without proof that names are unique.

- a) (4 pts) Let  $g : A \rightarrow C$  be an arrow in  $\mathcal{E}$ . Show that the subobject  $\langle \text{id}_A, g \rangle : A \rightarrow A \times C$  is named by the composite arrow  $\{\cdot\}g : A \rightarrow \Omega^C$ .
- b) (6 pts) Suppose that we have morphisms  $f : A \rightarrow B$ ,  $g : A \rightarrow C$  with  $f$  monic. Suppose  $h : B \rightarrow \Omega^C$  names the subobject  $\langle f, g \rangle : A \rightarrow B \times C$ . Prove that the square

$$\begin{array}{ccc} A & \xrightarrow{g} & C \\ f \downarrow & & \downarrow \{\cdot\} \\ B & \xrightarrow{h} & \Omega^C \end{array}$$

is a pullback. [Hint: in order to see that the diagram commutes, combine the square which testifies that  $h$  names  $\langle f, g \rangle$  with the pullback

$$\begin{array}{ccc} A & \xrightarrow{\text{id}_A} & A \\ \langle \text{id}_A, g \rangle \downarrow & & \downarrow \langle f, g \rangle \\ A \times C & \xrightarrow{f \times \text{id}_C} & B \times C \end{array}$$

and conclude that  $hf : A \rightarrow \Omega^C$  names the subobject  $\langle \text{id}_A, g \rangle : A \rightarrow A \times C$ .]

**Exercise 2** Recall that there is a map  $\wedge : \Omega \times \Omega \rightarrow \Omega$  which is such that if  $A$  and  $B$  are subobjects of  $X$ , classified by  $\phi, \psi : X \rightarrow \Omega$  respectively, then the intersection  $A \cap B$  is classified by the composite

$$X \xrightarrow{\langle \phi, \psi \rangle} \Omega \times \Omega \xrightarrow{\wedge} \Omega$$

Recall also that the order relation on  $\Omega$  is the subobject  $\Omega_1$  of  $\Omega \times \Omega$  which is the equalizer of  $\wedge$  and the first projection.

Prove that there is a monic arrow  $i_\Omega : \Omega \rightarrow \Omega_1$  which is a partial map classifier for  $\Omega$ .

**Exercise 3** a) (5 pts) Let  $\mathcal{C}$  be a small category and let  $\text{Cov}$  be a Grothendieck topology on  $\mathcal{C}$ . Suppose  $R \in \text{Cov}(C)$  contains a split epimorphism  $f : B \rightarrow C$ . Prove: for any presheaf  $X$  on  $\mathcal{C}$  and any compatible family  $(x_g \mid g \in R)$  indexed by  $R$  (so:  $x_g \in X(\text{dom}(g))$  for all  $g \in R$ ), there is a unique amalgamation in  $X(C)$ .

b) (5 pts) Use part a) to characterize all inclusions  $\mathcal{E} \rightarrow \widehat{\mathcal{G}}$ , where  $\mathcal{E}$  is a Grothendieck topos and  $\widehat{\mathcal{G}}$  is the topos of right  $\mathcal{G}$ -sets, for a group  $\mathcal{G}$ .

**Exercise 4** A *filter of subterminals* in a topos  $\mathcal{E}$  is a collection  $\Phi$  of subobjects of 1, satisfying the following conditions:

- i) if  $U \in \Phi$  and  $U \leq V$  then  $V \in \Phi$
- ii)  $1 \in \Phi$
- iii) if  $U, V \in \Phi$  then  $U \cap V \in \Phi$

a) (2 pts) Let  $\Phi$  be a filter of subterminals in a topos  $\mathcal{E}$ . For objects  $A, B$  of  $\mathcal{E}$  we define a  $\Phi$ -*morphism from  $A$  to  $B$*  to be an equivalence class of arrows  $A \times U \rightarrow B$  in  $\mathcal{E}$  with  $U \in \Phi$ ; two such arrows  $f : A \times U \rightarrow B$  and  $g : A \times V \rightarrow B$  are equivalent if there is a  $W \in \Phi$  with  $W \leq U \cap V$  and the square

$$\begin{array}{ccc} A \times W & \longrightarrow & A \times V \\ \downarrow & & \downarrow g \\ A \times U & \xrightarrow{f} & B \end{array}$$

commutes.

Show that there is a category  $\mathcal{E}_\Phi$  which has the same objects as  $\mathcal{E}$ , and  $\Phi$ -morphisms as arrows.

- b) (3 pts) Characterize the monos in  $\mathcal{E}_\Phi$ .
- c) (5 pts) Show that there is a functor  $P_\Phi : \mathcal{E} \rightarrow \mathcal{E}_\Phi$  which preserves finite limits and exponentials. [Hint: given a finite diagram of  $\Phi$ -morphisms, there is some  $U \in \Phi$  such that every  $\Phi$ -morphism in the diagram is equivalent to one of the form  $A \times U \rightarrow B$ ]

[Remark: in fact it can be shown that  $\mathcal{E}_\Phi$  is a topos and the functor  $P_\Phi$  is logical.]